

CS 61C  
Fall 2018RISC-V Control Flow  
Discussion 4: September 17, 2018

## 1 RISC-V with Arrays and Lists

Comment each snippet with what the snippet does. Assume that there is an array, `int arr[6] = {3, 1, 4, 1, 5, 9}`, which starts at memory address `0xBFFFFFFF00`, and a linked list struct (as defined below), `struct ll* 1st;`, whose first element is located at address `0xABCD0000`. `s0` then contains `arr`'s address, `0xBFFFFFFF00`, and `s1` contains `1st`'s address, `0xABCD0000`. You may assume integers and pointers are 4 bytes and that structs are tightly packed.

```
struct ll {
    int val;
    struct ll* next;
}
```

```
[1.1] lw t0, 0(s0)
      lw t1, 8(s0)
      add t2, t0, t1
      sw t2, 4(s0)
```

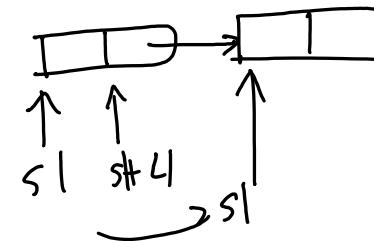
$$\left. \begin{array}{l} \text{arr}[0] \rightarrow t0 \\ \text{arr}[2] \rightarrow t1 \\ t0+t1 \rightarrow t2 \\ t2 \rightarrow \text{arr}[1] \end{array} \right\} \rightarrow \text{arr}[1] = \text{arr}[0] + \text{arr}[2]$$

if ( $s1 == 0$ ) → jump to end

```
[1.2] loop: beq s1, x0, end
        lw t0, 0(s1)
        addi t0, t0, 1
        sw t0, 0(s1)
        lw s1, 4(s1)
        jal x0, loop
end:
```

$\text{MEM}(s1) \rightarrow t0$   
 $t0+1 \rightarrow t0$   
 $t0 \rightarrow \text{MEM}(s1)$   
 $s1 = s1 \rightarrow \text{next}$

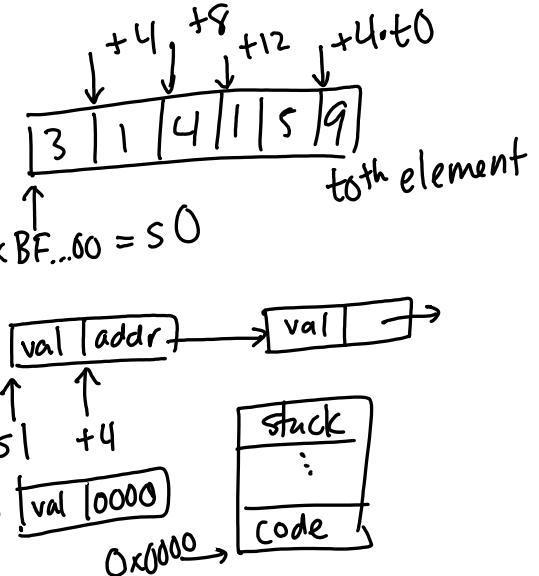
$\Rightarrow$  incr. all elements of `ll` by 1



```
[1.3] loop: slti t1, t0, 6
      beq t1, x0, end
      slli t2, t0, 2
      add t3, s0, t2
      lw t4, 0(t3)
      sub t4, x0, t4
      sw t4, 0(t3)
      addi t0, t0, 1
      jal x0, loop
end:
```

0 → t0  
if  $t0 < 6 \Rightarrow 1 \rightarrow t1$ , else 0 → t1  
 $t1 == 0?$   
4 · t0 → t2  
 $t2 + s0 \rightarrow t3$   
 $\text{MEM}(t3) \rightarrow t4$  ← get 6<sup>th</sup> element  
0 - t4 → t4  
 $t4 \rightarrow \text{MEM}(t3)$  ← put it back  
negate 6 elements

lw t4, 0 (  $s0+t2$  )  
 $\underbrace{\phantom{0}}_{t3}$



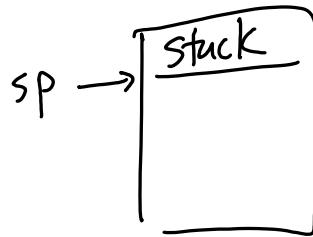
## 2 RISC-V Calling Conventions

- 2.1 How do we pass arguments into functions?

arg registers: a0 - a7

- 2.2 How are values returned by functions?

a0, a1



- 2.3 What is sp and how should it be used in the context of RISC-V functions?

points to where we can save values

- 2.4 Which values need to be saved by the caller, before jumping to a function using jal?

t0-t6, a0-a7, (ra)

- 2.5 Which values need to be restored by the callee, before using jalr to return from a function?

Save s0-s11, sp, gp, tp, (ra)

## 3 Writing RISC-V Functions

- 3.1 Write a function sumSquare in RISC-V that, when given an integer n, returns the summation below. If n is not positive, then the function returns 0.

$$n^2 + (n-1)^2 + (n-2)^2 + \dots + 1^2$$

For this problem, you are given a RISC-V function called square that takes in an integer and returns its square. Implement sumSquare using square as a subroutine.

sumSquare:

|      |             |   |          |
|------|-------------|---|----------|
| addi | sp, sp, -12 | ] | prologue |
| sw   | s0, 0(sp)   |   |          |
| sw   | s1, 4(sp)   |   |          |
| sw   | ra, 8(sp)   |   |          |

jal ra, square

end:

|      |            |   |          |
|------|------------|---|----------|
| add  | a0, s1, x0 | ] | epilogue |
| lw   | s0, 0(sp)  |   |          |
| lw   | s1, 4(sp)  |   |          |
| lw   | ra, 8(sp)  |   |          |
| addi | sp, sp, 12 |   |          |

jr ra

s0, s1, ra

body

return value

## 4 More Translating between C and RISC-V

- 4.1 Translate between the C and RISC-V code. You may want to use the RISC-V Green Card as a reference. We show you how the different variables map to registers – you don't have to worry about the stack or any memory-related issues.

| C   | RISC-V |
|---|--------|
| <pre>// Nth_Fibonacci(n): // s0 -&gt; n, s1 -&gt; fib // t0 -&gt; i, t1 -&gt; j // Assume fib, i, j init'd to: int fib = 1, i = 1, j = 1; if (n==0)     return 0; else if (n==1)     return 1; n -= 2; while (n != 0) {     fib = i + j;     j = i;     i = fib;     n--; } return fib;</pre> |        |